

“Hello World” Program

'data' and 'text'
areas

```
## entry point of program must be called 'main'
main: addiu $v0, $0, 4      # syscall 4 (print_str)
      la $a0 msg
      syscall
      addiu $v0, $0, 10
      syscall              # syscall 10 (exit)
```

Syscalls -- used for I/O to console, exit.

\$v0 holds syscall
identification number

Argument registers used to pass values to syscall

'syscall' is like a software interrupt

Service	Call Code	Arguments	Result
print_int	1	\$a0=integer	
print_float	2	\$f12=float	
print_double	3	\$f12=double	
print_string	4	\$a0=string	
read_int	5		int, (in \$v0)
read_float	6		float (in \$f0)
read_double	7		double (in \$f0)
read_string	8	\$a0=buffer, \$a1=length	
sbrk	9	\$a0 = amount	addr (in \$v0)
exit	10		

la pseudoinstruction

la – **load address**, is an assembler pseudoinstruction

Useful for placing the address of an operand into a register:

```
la $a0 msg
```

Actually translates into:

```
lui $1, upper_16_bits_msg_addr
```

```
ori $a0, $1, lower_16_bits_msg_addr
```

where value represents the starting address of the 'msg' operand.
If lower half of address is '0000h', then the 'ori' instruction is not generated.

Note that Register '1' is used by the pseudoinstruction !!!!

Register Definitions

Name	Reg Number	Usage	Preserved on call?
\$zero	0	constant value 0	N.A.
\$v0-\$v1	2-3	values for results, expression eval	no
\$a0-\$a3	4-7	arguments	yes
\$t0-\$t7	8-15	temporaries	no
\$s0-\$s7	16-23	saved	yes
\$t8-\$t9	24-25	more temps	no
\$gp	28	global pointer	yes
\$sp	29	stack pointer	yes
\$fp	30	frame pointer	yes
\$ra	31	rtm address	yes

```
.data
msg1: .asciiz "Enter a string: \n "    Echo string
buff: .space 64                       program

.text
## entry point of program must be called 'main'
main: addiu $v0, $0, 4    # syscall 4 (print_str)
      la $a0 msg1        # prompt user
      syscall
      addiu $v0, $0, 4

      la $a0 buff        #get string from user
      addiu $a1, $0, 64
      addiu $v0, $0, 8    #syscall 8 (read str)
      syscall

# echo string
      addiu $v0, $0, 4
      syscall

#exit
      addiu $v0, $0, 10
      syscall                # syscall 10 (exit)
```

Implementing a Stack

SPIM initialization sets SP to high memory (\$7FFFeffc)

To push a register value onto the stack:

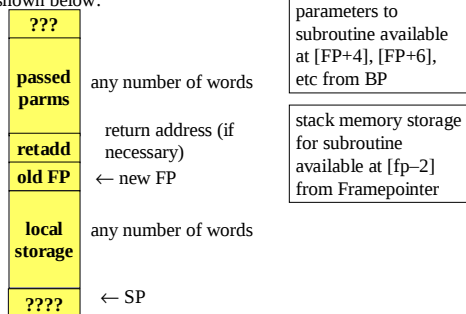
```
addiu $sp,$sp,-4    # decrement stack pointer
sw    $6, 0($sp)    # push register 6 on stack
```

To pop a value off the stack

```
lw    $6, 0($sp)
addiu $sp, $sp, 4    # increment stack pointer
```

A Stack Frame

A generalized picture of a *stack frame* used by a subroutine is shown below:



Code for Stack Frame (assume 2 words on stack as passed parameters)

```
suba:
    addiu $sp,$sp, -4    ← Save RTN address (if needed)
    sw    $ra, 0($sp)
    addiu $sp,$sp, -4    ← Save old framepointer
    sw    $fp, 0($sp)
    addiu $fp,$sp,$0     ← Get new framepointer
    addiu $sp, $sp,-8    ← Allocate local storage (if needed)
    ...
    addiu $sp,$fp, 0     ← Restore SP, free local storage
    lw    $fp, 0($sp)    ← restore old framepointer
    addiu $sp,$sp, 4     ← pop return address and clean
    lw    $ra, 0($sp)    stack of passed parameters
    addiu $sp, $sp, 12
    $jr   $ra            ← return
```

High Level Support for Stack Frames

- Some processors have special instructions that support stack frames

- x286 processors and higher has 'Enter' and 'Leave' instructions

<pre>suba proc push bp mov bp, sp sub sp, 8 mov sp, bp pop bp ret 6</pre>	 	<pre>suba proc enter 8, 0 leave ret 6</pre>
---	--	---

Not the only possible arrangement of stack frames

- Optimizing compilers will take short cuts on stack frames
- Will leave out parts of frame if not needed
 - ⇒ Saving of return address
 - ⇒ Framepoint
 - ⇒ Local Variable storage
 - ⇒ Passing of arguments on stack
- One or more of these features can be left out
